# How to Cross Open Spaces? 

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## Introduction

Finding the shortest path in open space is a well known challenge for pedestrian routing engines [1, 2]. A common solution is routing on the bounded polygon edges, which causes in most cases an unnecessarily long route (figure 1a). A possible solution is to create a subgraph crossing the open space. This research project assesses this approach and investigates its implications for routing engines.

## Method

There are many algorithms for creating subgraphs [1]. Figure 1 shows some examples. They were created with JTS Topology Suite. The dotted black line shows the shortest route from the start to the destination while the graph is visualized by the thin black lines.

## Algorithms

Delaunay, Voronoi and Visibility All take all nodes of the polygon into account. In contrast, Visibility GCP uses just points which are connected to other parts of the graph (Graph Connection Points). Spider is an extension of Grid with additional diagonals [2]. Table 1 shows that the standard route is $29 \%$ longer than the direct route with the Visibility All algorithm.

(a) Standard

(d) Grid 5 meter

(g) Spider 5 meter

(b) Voronoi

(e) Grid 10 meter

(h) Spider 10 meter

(c) Delaunay

(f) Visibility GCP

(i) Visibility All

Figure 1: Open space algorithms.
Table 1: Lengths of the routes and difference to Visibility All. Additional graph edges compared to Standard.

| Algorithm | Length | Graph Edges |
| :--- | :--- | ---: |
| Visibility All | $161 \mathrm{~m}(+0 \%)$ | +27 |
| Visibility GCP | $172 \mathrm{~m}(+7 \%)$ | +9 |
| Delaunay | $169 \mathrm{~m}(+5 \%)$ | +8 |
| Grid 5 meter | $197 \mathrm{~m}(+22 \%)$ | +504 |
| Grid 10 meter | $197 \mathrm{~m}(+22 \%)$ | +131 |
| Standard | $208 \mathrm{~m}(+29 \%)$ | +0 |
| Voronoi | $197 \mathrm{~m}(+22 \%)$ | +19 |
| Spider 5 meter | $168 \mathrm{~m}(+4 \%)$ | +1034 |
| Spider 10 meter $168 \mathrm{~m}(+4 \%)$ | +277 |  |

## Analysis

In order to compare these algorithms on a larger scale, we implemented them in the GraphHopper routing engine. Afterwards we computed the routing graph for each algorithm. The test area was the OpenStreetMap dataset of BadenWürttemberg and Austria. The processing time and the count of additional edges were compared with the respective values of the standard graph.

However, only the polygons stored as ways (closed linestrings) were considered. Polygons that are stored as relations (a data type which consists of a set of OpenStreetMap features) were not considered due to technical difficulties. Nevertheless, the analysis is still representative, since most of the open spaces are stored as ways (table 2).


Figure 2 : Routing possibilities through the open space.

## Results

The diagrams (figure 3) show that both the edge count and the computation time increases for most of the algorithms. The additional edge count of the algorithms shows a similar pattern for both regions. In contrast, the additional computation time for the Visibility All is much higher in Austria than in Baden-Württemberg.

Table 2 shows that there are less open spaces in Austria than in Baden-Württemberg. This explains why the relative increase of the edge count in Austria is lower. The quality of the routing result can be assessed with figure 1 and table 1. The Visibility All algorithm generates the most natural route and is moreover the shortest one. However, its computation time might take very long (figure 3 (b)).

Table 2 : Facts about test areas.

|  | Austria | Baden-Württemberg |
| :--- | :---: | :---: |
| standard comp. time | 8 min 51 s | 13 min 50 s |
| graph edges | 2.5 million | 2.7 million |
| open spaces total | 2693 | 5641 |
| open spaces ways | 2459 | 5225 |
| open spaces relations | 234 | 416 |
| highway ways | 1.44 million | 1.46 million |
| share open spaces | $0.19 \%$ | $0.38 \%$ |


(a) Baden-Würtemberg

(b) Austria

Figure 3 : Processing results (sorted by edge count).

## Outlook

This work has done a pre-assessment for implementing algorithms into a routing engine. The Visibility All algorithm looks most promising so far. However, it should be improved by removing unnecessary connections [1]. This would decrease the additional edge count and possibly the computation time. In a next step, areas modeled as relations should be taken into account as well. Moreover it should be tested if open space routing has an effect on the performance of routing requests.

## References

[1] Graser, Anita (2016). Integrating Open Spaces into OpenStreetMap Routing Graphs for Realistic Crossing Behaviour in Pedestrian Navigation, Gl_Forum Journal for Geographic Information Science.
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## 1

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